Shifting Segments to Optimality

Stefan Felsner

Abstract We begin with two problems which do not appear to be related. Then we use the 'air-pressure' method to prove a theorem about rectangular dissections with prescribed rectangle areas. A corollary to this theorem characterizes area-universal rectangular dissections. These dissections happen to be central to the solution of the two problems.

1 Introduction and Two Problems

A rectangular dissection is a partition of a frame rectangle into rectangles, Fig. 1 shows an example. Rectangular dissections are studied in various fields, see Fig. 2.

- Architects look at them in the context of floorplan generation [11, 14].
- Floorplaning is relevant for module placement in VLSI design [4, 21].
- In graph drawing, rectangular dissections play a role in various representation models for planar graphs [8, 12].
- In cartography, rectilinear dissections are studied as a special class of cartograms [13, 17, 19].

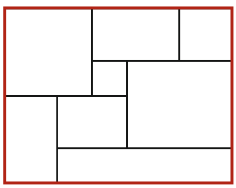

Fig. 1 A rectangular dissection

S. Felsner (✉)
Institut für Mathematik, Technische Universität Berlin, Straße des 17. Juni 136, 10623 Berlin, Germany
e-mail: felsner@math.tu-berlin.de

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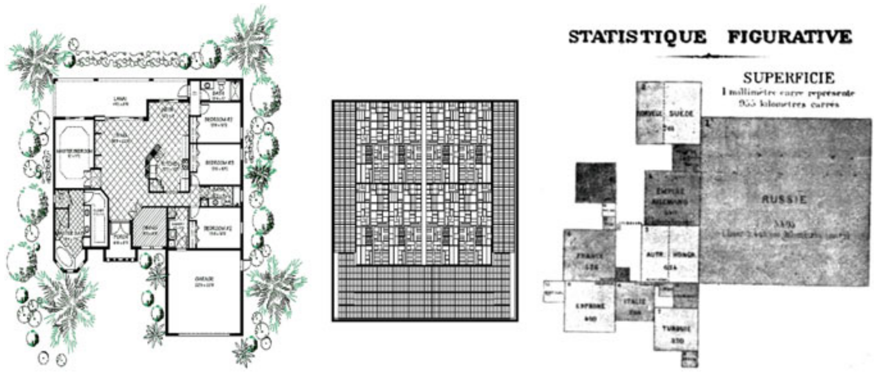

Fig. 2 Near rectangular dissections in applications

In the applications the areas of the rectangles of a dissection are relevant. In many cases these areas are prescribed. A rectangular dissection is *area-universal* if for any assignment of positive weights to the rectangles there is a combinatorially equivalent dissection such that the areas of rectangles are equal to the given weights.

Central to this chapter is the characterization of area-universal rectangular dissections (Theorem 5). In Sect. 2 we state the theorem and discuss some proofs and generalizations. Before getting there we present two problems which do not appear to have much in common. In Sect. 3 we show that both problems can be solved by clever applications of the theorem.

First Problem

By default a *dissection* shall be a rectangular dissection. A dissection is *generic* if it has no cross, i.e., no point where four rectangles of the partition meet. A *segment* of a dissection is a maximal nondegenerate interval that belongs to the union of the boundaries of the rectangles. In general we disregard the four segments from the boundary frame, i.e., we only consider inner segments. Segments are either hori-

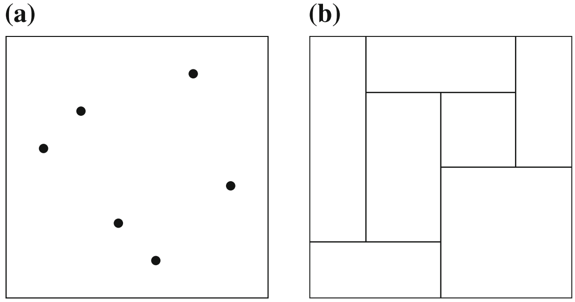

Fig. 3 A generic set of six points and a generic dissection with six segments

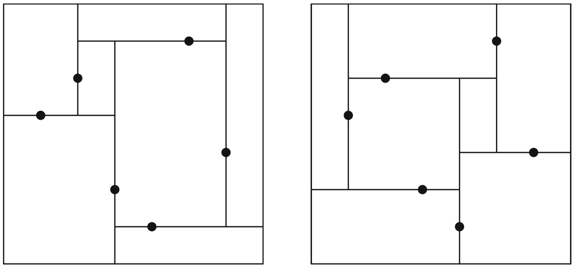

Fig. 4 Two cover maps from the dissection of Fig. 3b to the point set of Fig. 3a

zontal or vertical. The segments of a generic dissection are internally disjoint. Two dissections R and R' are *weakly equivalent* if there exists a bijection ϕ between their segments that preserves the orientation (horizontal/vertical) and such that a segment s has an endpoint on a segment t in R iff $\phi(s)$ has an endpoint on $\phi(t)$ in R'. A set P of points in $\mathbb{R}^2$ is *generic* if no two points from P have the same x or y coordinate, see Fig. 3.

Let P be a set of n points in a rectangular frame F and let R be a generic dissection with n segments. A *cover map* from R to P is a dissection R' that is weakly equivalent to R and has outer rectangle F such that every segment of R' contains exactly one point from P. Figure 4 shows an example.

Problem 1 Does a cover map from R to P exist for all pairs (R, P) where R is a generic dissection with n segments and P is a generic set of n points?

Second Problem

This problem is about *rectilinear duals* of planar graphs. In this drawing model the vertices are represented by simple rectilinear polygons, while edges are represented by side-contacts between the corresponding polygons, see Fig. 5.

Now assume that positive weights $w(v)$ have been assigned to the vertices of the graph. A *rectilinear cartogram* is a rectilinear dual with the additional property that for all vertices the area of the polygon representing v equals $w(v)$. A relevant parameter measuring the complexity of a cartogram is the maximum number of sides of any polygon.

Problem 2 What is the minimum number k such that any given planar triangulation, i.e., maximally planar graph, with positive weights $w(v)$ admits a rectilinear cartogram with $\leq k$-sided polygons in a rectangular frame F of area $\sum_v w(v)$?

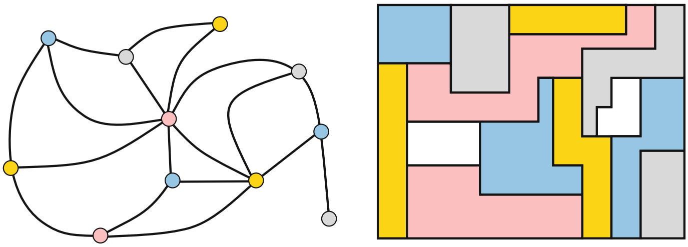

Fig. 5 A graph with a rectilinear dual containing two holes, the *white* regions

2 Area-Universality of Weak Equivalence Classes

The following theorem is a formalization of the statement given in the heading: *weak equivalence classes are area-universal.*

Theorem 3 *Let R be a dissection with rectangles $r_1, \dots, r_{n+1}$, let the frame F be a rectangle and let $w : \{1, \dots, n+1\} \to \mathbb{R}_+$ be a weight function with $\sum_i w(i) = \mathsf{area}(F)$. There exists a unique dissection R' contained in F that is weakly equivalent to R such that the area of the rectangle r'_i in R' is $w(i)$.*

We will discuss a proof based on an iterative approach below. Before getting there, however, we introduce a class of dissections so that we can state the most important special case of the theorem.

Two dissections are *dual equivalent* if they have the same dual graph. In most applications we are interested in finding an appropriate member of the dual equivalence class. If two dissections are weakly equivalent they need not be dual equivalent, for example in Fig. 4 the rectangle in the lower left corner has 4 neighbors in the left dissection but only 3 in the right dissection.

A segment s of a dissection is *one-sided* if s is the side of at least one of the rectangles, in other words all the segments that have an endpoint on s are on the same side of s. A dissection is *one-sided* if every segment of the dissection is one-sided. The following observation was made in [6].

Proposition 4 *All dissections in the weak equivalence class of a one-sided dissection are dual equivalent.*

Together with Theorem 3 this yields the key theorem.

Theorem 5 *One-sided dissections are area-universal.*

With the following definitions we set the stage for a generalization of Theorem 3. Let $\mu : [0, 1]^2 \to \mathbb{R}_+$ be a density function on the unit square with mass 1, i.e.,

$\int_0^1 \int_0^1 \mu(x, y) dx dy = 1$. We assume that μ is well behaved so that all the integrals we need exist and are positive. The *mass* of an axis aligned rectangle $r \subseteq [0, 1]^2$ is defined as $m(r) = \iint_r \mu(x, y) dx dy$.

Theorem 6 *Let $\mu : [0, 1]^2 \to \mathbb{R}_+$ be a density function on the unit square. If R is a dissection with rectangles $r_1, \ldots, r_{n+1}$ and $w : \{1, \ldots, n+1\} \to \mathbb{R}_+$ a positive weight function with $\sum_{i=1}^{n+1} w(i) = 1$, then there exists a unique dissection R' in the unit square that is weakly equivalent to R such that the mass of the rectangle r_i' in R' is exactly $w(i)$.*

A full proof of the theorem is given in [9]. Here we only sketch the proof, it is based on a force directed method that exploits a physical analogy with air-pressure. Consider a representation of R in the unit square and compare the mass $m(r_i)$ to the intended mass $w(i)$. The quotient of these two values can be interpreted as the pressure inside the rectangle. Integrating this pressure along a side of the rectangle yields the force by which r_i is pushing against the segment that contains the side. The difference of pushing forces from both sides of a segment yields the effective force acting on the segment. The intuition is that shifting a segment in the direction of the effective force yields a better balance of pressure in the rectangles. We show that iterating such improvement steps drives the realization of R towards a situation with $m(r_i) = w(i)$ for all i, i.e., the procedure converges towards the dissection R' whose existence we want to show.

Let $r_i = [x_l, x_r] \times [y_b, y_t]$ be a rectangle of R. The pressure $p(i)$ in r_i is $p(i) = \frac{w(i)}{m(r_i)}$. Let s be a segment of R and let r_i be one of the rectangles with a side in s. Let s be vertical with x-coordinate x_s and let $s \cap r_i$ span the interval $[y_b(i), y_t(i)]$. The (undirected) *force $f(s, i)$ imposed on s by r_i* is the pressure $p(i)$ of r_i times the density dependent length of the intersection, i.e.,

$$f(s, i) = p(i) \int_{y_b(i)}^{y_t(i)} \mu_{x_s}(y) dy.$$

The *force $f(s)$ acting on s* is obtained as a sum of the directed forces imposed on s by incident rectangles.

$$f(s) = \sum_{r_i \text{ left of } s} f(s, i) - \sum_{r_i \text{ right of } s} f(s, i).$$

Symmetric definitions apply to horizontal segments.

Balance for Rectangles and Segments

A segment s is in *balance* if $f(s) = 0$. A rectangle r_i is in *balance* if $m(r_i) = w(i)$, i.e., if $p(i) = 1$.

Lemma 7 *All rectangles r_i of R are in balance if and only if all segments are in balance.*

Proof We only show one direction. Since all rectangles are in balance we can eliminate the pressures from the definition of the $f(s, i)$. With this simplification we get for a vertical segment s

$$f(s) = \sum_{r_i \text{ left of } s} \int_{y_b(i)}^{y_t(i)} \mu_{x_s}(y)dy - \sum_{r_j \text{ right of } s} \int_{y_b(j)}^{y_t(j)} \mu_{x_s}(y)dy.$$

Hence $f(s) = M_s - M_s = 0$, where M_s is the integral of the fiber density μ_{x_s} along s.

□

Balancing Segments and Optimizing the Entropy

Proposition 8 *If a segment s of R is unbalanced, then we can keep all the other segments at their position and shift s parallel to a position where it is in balance. The resulting dissection R' is weakly equivalent to R.*

The *entropy* of a rectangle r_i of R is defined as $-w(i) \log p(i)$. The *entropy* of the dissection R is

$$E = \sum_i -w(i) \log p(i)$$

The proof of Theorem 6 is in five steps:

(1) The entropy E is always nonpositive.
(2) $E = 0$ if and only if all rectangles r_i of R are in balance.
(3) Shifting an unbalanced segment s into its balance position increases the entropy.
(4) The process of repeatedly shifting unbalanced segments into their balance position makes R converge to a dissection R' such that the entropy of R' is zero.
(5) The solution is unique.

3 Solutions to the Problems

Mapping Segments on Points

Let R be a generic dissection with n segments. Let P be a generic set of n points in a rectangle F. Recall that a *cover map* from R to P is a dissection R' with outer rectangle F that is weakly equivalent to R such that every segment of R' contains exactly one point from P. The following theorem from [9] answers our first problem.

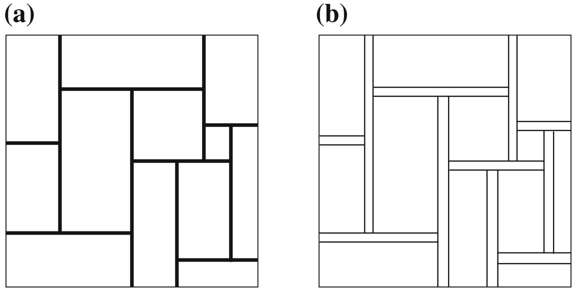

Fig. 6 Dissections R and the dissection R_S obtained by doubling the segments

Theorem 9 *If R is a generic dissection with n segments and P is a generic set of n points in a rectangle F, then there is a cover map from R to P.*

To be able to use Theorem 6 we first transform the point set P into a suitable density distribution $\mu = \mu_P$ inside F. This density is defined as the sum of a uniform distribution μ_1 with $\mu_1(q) = 1/\text{area}(F)$ for all $q \in F$ and a distribution μ_2 that represents the points of P. Choose some $\Delta > 0$ such that for all $p, p' \in P$ we have $|x_p - x_{p'}| > 3\Delta$ and $|y_p - y_{p'}| > 3\Delta$, this is possible because P is generic. Define $\mu_2 = \sum_{p \in P} \mu_p$ where $\mu_p(q)$ takes the value $(\Delta^2\pi)^{-1}$ on the disk $D_\Delta(p)$ of radius Δ around p and value 0 for q outside of this disk.

For a density ν over F and a rectangle $r \subseteq F$ we let $\nu(r)$ be the integral of the density ν over r. Using this notation we can write $\mu_1(F) = 1$ and $\mu_p(F) = 1$ for all $p \in P$, hence the total mass of F is $\mu(F) = 1 + n$.

Next we transform the dissection R into a dissection R_S. To this end we replace every segment by a thin rectangle, see Fig. 6. Let $\mathscr{S}$ be the set of new rectangles.

Define weights for the rectangles of R_S as follows. If R_S has r rectangles we define $w(r) = 1 + 1/r$ if $r \in \mathscr{S}$ and $w(r) = 1/r$ for all the rectangles of R_S that came from rectangles of R. Note that the total weight, $\sum_r w(r) = 1 + n$, is in correspondence to the total mass $\mu(R)$.

The data R with μ and R_S with w constitute, up to scaling of R and w, a set of inputs for Theorem 6. From the conclusion of the theorem we obtain a dissection R'_S weakly equivalent to R_S such that $m(r) = \iint_r \mu(x, y)dxdy = w(r)$ for all rectangles r of R'_S.

The definition of the weight function w and the density μ is so that R'_S should be close to a cover map from R to P: Only the rectangles $r \in \mathscr{S}$ that have been constructed by inflating segments may contain a disk $D_\Delta(p)$ and each of these rectangles may contain at most one of the disks. This suggests a correspondence $\mathscr{S} \leftrightarrow P$. However, a rectangle $r \in \mathscr{S}$ can use parts of several discs to accumulate mass. To find a correspondence between $\mathscr{S}$ and P we define a bipartite graph G whose vertices are the points in P and the rectangles in $\mathscr{S}$:

- A pair (p, r) is an edge of G iff $r \cap D_\Delta(p) \neq \emptyset$ in R'_S.

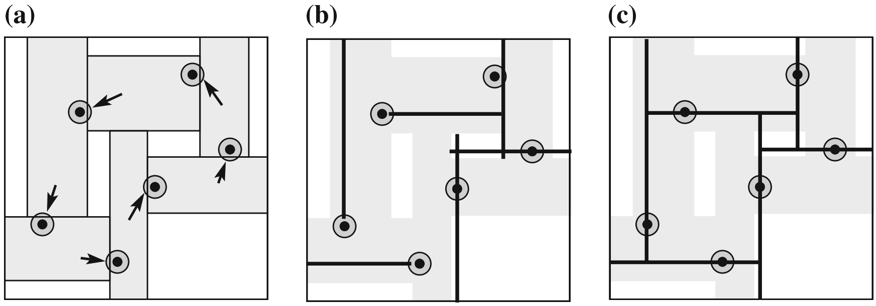

Fig. 7 (**a**) A solution R'_S with a matching indicated by the arrows. (**b**) Segments are shifted to the corresponding points. (**c**) Small final adjustments (clipping and enlarging) yield R'.

The proof of the theorem is completed by proving two claims:

- G admits a perfect matching.
- From R'_S and a perfect matching M in G we can produce a dissection R' that realizes the cover map from R to P.

For the first of the claims we check Hall's matching condition. Consider a subset A of $\mathscr{S}$. Since R_S realizes the prescribed weights we have $m(A) = \mu(A) = \sum_{r \in A} \mu(r) = \sum_{r \in A} w(r) = |A|(1 + 1/r)$. Since $\mu_1(A) < 1$ and $\mu_p(A) \leq 1$ for all $p \in P$, there must be at least $|A|$ points $q \in P$ with $\mu_q(A) > 0$, these are the points that have an edge to a rectangle from A in G. We have thus shown that every $A \subset \mathscr{S}$ is incident to at least $|A|$ points in G, hence, there is an injective mapping $\alpha : \mathscr{S} \to P$ such that $r \cap D_\Delta(\alpha(r)) \neq \emptyset$ in R'_S for all $r \in \mathscr{S}$.

Given the matching α the construction of the dissection R' that realizes the cover map from R to P is completed in two further steps, see Fig. 7b, c.

Cartograms with Optimal Complexity

In a series of papers the complexity of polygons used for the cartograms of triangulations has been reduced from 40 to 34 then 12 and 10. Finally, in [3] the following optimal result was obtained.

Theorem 10 *Every planar triangulation admits an area-universal rectilinear cartogram with* $\leq$*8-sided polygons.*

The construction is fairly easy with the right tools at hand. First we need a Schnyder wood of the input triangulation G. Let a_1, a_2, a_3 be the outer vertices of G, an orientation and coloring of the inner edges with 3 colors (we identify colors (1, 2, 3) with (red, green, blue)) is a *Schnyder wood* if:

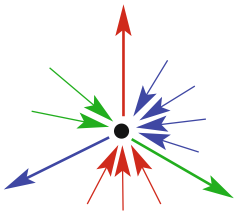

Fig. 8 Schnyder's edge coloring rule

(1) All edges incident to an outer vertex a_i are in-edges and colored i.
(2) Every inner vertex v has three outgoing edges colored red, green and blue in clockwise order. All the incoming edges in an interval between two outgoing edges are colored with the third color, see Fig. 8.

These structures were defined by Schnyder in [15], where it was also shown that every triangulation admits a Schnyder wood. Moreover, if T_i is the set of oriented edges of color i and T_i^{-1} is the same set with reversed orientations, then it holds that $T_1 \cup T_2^{-1} \cup T_3^{-1}$ is acyclic. This property can be used to show that every triangulation has a contact representation with internally disjoint $\perp$ shapes. Figure 9 shows an example.

The $\perp$-representation can be viewed as a rectangular dissection. Now replace every segment of this dissection by a thin rectangle. This yields a one-sided dissection R_G, see Fig. 10(left). With a vertex v of G we associate the polygon P_v formed as the union of four rectangles. These are the two rectangles that were obtained from the two segments of the $\perp$ shape representing v together with the two rectangles that have parts of the horizontal segment of this $\perp$ as bottom side. In Fig. 10(right) the polygons P_v are distinguished by colors.

It is easily checked that the polygons P_v have at most 8 corners, hence, at most 8 sides. Given a set of weights $w : V \to \mathbb{R}^+$ we can arbitrarily break $w(v)$ into four positive values and assign these to the rectangles constituting P_v. Since the dissection R_G is one-sided and, hence, area-universal there is a realization of the dissection where the area of $P(v)$ equals $w(v)$.

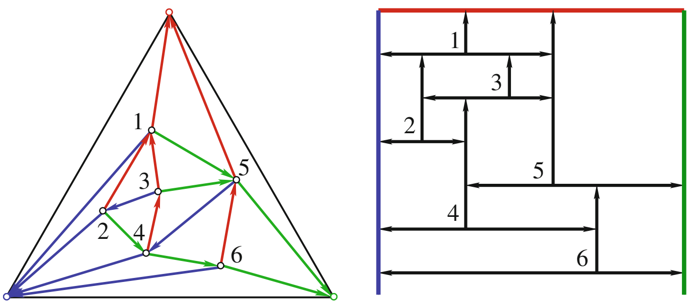

Fig. 9 A triangulation with a Schnyder wood and a $\perp$-representation

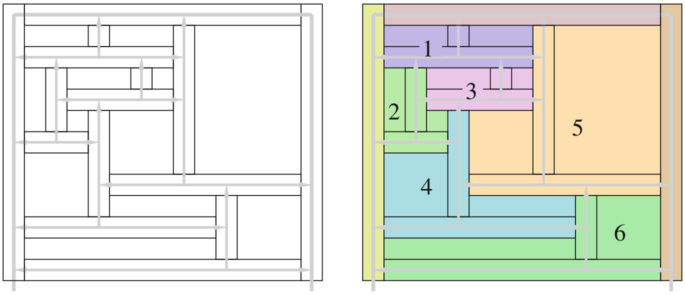

Fig. 10 The one-sided dissection resulting from the ⊥-representation of Fig. 9

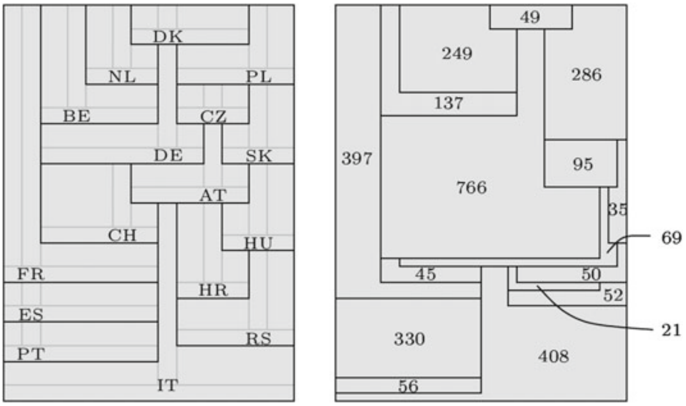

Fig. 11 Central European states represented by polygons with at most 8 sides (*left*). A cartogram where the areas are proportional to the emission of CO_2 in 2009 (*right*).

From the thesis of Torsten Ueckerdt [20] we borrow Fig. 11, which shows a cartogram displaying real data. The cartogram was computed with the method of this section.

4 Background and Additional Problems

Theorem 3 was first proven by Wimer et al. [21]. They take the width x_i and height y_i of the rectangle r_i as variables and show that the system consisting of linear equations which correspond to left-to-right and bottom-to-top sequences of rectangles together with the non-linear equations $x_i y_i = w(i)$ has a unique solution. The theorem was rediscovered by Eppstein et al. [5]. They prove it with an argument based on "*invariance of domain*." Both proofs are purely existential. However, in [5, 6] it is noted that the solution can be computed by iteratively reducing the distance between the present weights and the intended weight vector $(w(i))_i$. The 'air-pressure' method proposed by Izumi, Takahashi and Kajitani [10] is such an iterative approach. In [9] the air-pressure technique was used to prove Theorem 6. The sketch given here is based on this paper. A short non-constructive proof of Theorem 6 was given by Schrenzenmaier [16, p. 21], he adapted the proof of Theorem 3 from [5].

The two main problems regarding area-universal rectangular dissections are the following:

- Given R and w, is it possible to compute the weakly equivalent dissection R' realizing the weights in polynomial time (efficient Theorem 3)?
- Characterize graphs that admit a one-sided dual or find a polynomial recognition algorithm for them.

Beside this it would be very interesting to identify further instances of area-universality. Two such instances are straight line drawings of 3-regular planar graphs with prescribed face areas (Thomassen [18]) and straight line drawings of grids with prescribed face areas, a.k.a. table cartograms (Evan et al. [7]).

Problem 1 was a conjecture of Ackerman et al. [1]. They were motivated by the study of the function $Z(P)$ counting rectangulations of a generic point set P. Combining results from [9] (lower) and [1] (upper) we know that $Z(P)$ is in $\Omega(8^{n+1}/(n+1)^4)$ and in $O(20^n/n^4)$. The lower bound is tight for some sets P, to improve the upper bound remains an intriguing problem.

The construction of area-universal rectilinear cartograms with $\leq$8-sided polygons is from Alam et al. [3]. As already noted the construction based on our key theorem is not known to be efficient. Polynomial constructions of cartograms with $\leq$8-sided polygons are known for Hamiltonian triangulations [3] and with $\leq$10-sided polygons for general triangulations [2]. Is there a polynomial algorithm for constructing cartograms with $\leq$8-sided polygons for general triangulations?

Recognition of planar graphs which admit rectangular cartograms or cartograms with $\leq$6-sided polygons is also wide open.

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