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Corrections to “Introduction to Cryptography, Second Edition”

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In the line before Example 1.7.4, replace $a_i$ by $\alpha_i$.

Proof of Theorem 2.9.5: Theorem 2.9.2 instead of Theorem 2.9.3 (Twice)

Proof of Corollary 2.11.3: Theorem 2.9.2 instead of Theorem 2.9.3.

Lemma 2.19.2: Use a “plain” $K$.

Exercise 2.22.12: $d_i$ is missing in the sum.

At the bottom the sequence reads $c_1, c_2, \ldots, c_n$. The last entry should be $c_u$ instead.

Equation (3.3): replace $z_{i-j}$ by $s_{i-j}$.

above example 3.9.3, the $p_i$ should be $c_i$.

line 2 of 3.13: The name is Blaise de Vigenère.

Example 3.14.1: The determinant of $A$ is even, and so the cipher is not allowable since it is not relatively prime to $m = 26$. Replace FUSS replaced by FOOT.

line 2: Pr($a$) instead of $P(a)$.

p. 105, line 2 of Definition 4.2.2: The “end quote” should be placed after ‘occurs’ (and not after the $B$).

line 1 of Example 4.2.3: Delete “probability of the”.

line 9 from below: $m$ should be replaced by $p$ (3 times).

Figure 5.1: replace “Expansionsfunktion” by “expansion function”, “S-Boxen” by “S-boxes” and $f(R, K)$ by $f(K, R)$.

In Table 5.3, description of the function $P$ the positions for 10 and 20 must be switched.

Replace $f(R_0, K_1)$ by $f(K_1, R_0)$.

4th last and 2nd last lines of Section 5.3: In both strings, the 3rd and 16th bits (from the left) should be changed (that’s a result of the problem with the P-table).

Those arrays have “four” rows …x

line 3 of Example 7.2.1: Read $\gcd(3, 220) = 1$. p. 145, line 3 of Example 7.2.5:

Example 8.3.5: 119 should be replaced by 110 (twice), and 26 by 165.

line 8 of 2nd paragraph: $1024$ instead of $1023$.

line 7 of Section 8.5.4: Read $K = A^b \mod p$.

last line of first paragraph: Read $g^c \equiv g^{ab} \mod p$.

line 1: $b \in \{0, 1, \ldots, p - 2\}$.

In equation (10.4) $a +$ is missing:

$$p^{c-1}x = x_0p^{c-1} + p^c(x_1 + x_2p + \ldots + x_{c-1}p^{c-2}).$$

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$$p^{c-1}x = x_0p^{c-1} + p^c(x_1 + x_2p + \ldots x_{c-1}p^{c-2}).$$

Line 15: In the definition of SHA-1 we have

$$C = S^{30}(B)$$
instead of

\[ C = S^{36}(B). \]

**p. 279** Exercise 12.9.5: In the ElGamal signature scheme use the prime number \( p \) and the primitive root \( g \mod p \). Suppose that \( p \equiv 1 \pmod{4} \) and that \( g \) has only small prime factors. Let \( A \) be Alice's public key.

1. Show that a solution \( z \) of the congruence \( A^q = g^{qz} \mod p \) can be found efficiently.
2. Let \( x \) be a document and let \( h \) be its hash value. Prove that \( (q, (p - 3)(h - qz)/2) \) is a valid signature of \( x \).

**p. 295** The correct formula for the determinant of the Vandermonde matrix is

\[
\det U = \prod_{1 \leq i < j \leq \ell} (x_j - x_i).
\]
Introduction to Cryptography
Buchmann, J.
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